

General Certificate of Education (A-level) June 2012

Computing
COMP3
(Specification 2510)
Unit 3: Problem Solving, Programming,
Operating Systems, Databases and Networking

## Final

Mark schemes are prepared by the Principal Examiner and considered, together with the relevant questions, by a panel of subject teachers. This mark scheme includes any amendments made at the standardisation events which all examiners participate in and is the scheme which was used by them in this examination. The standardisation process ensures that the mark scheme covers the candidates' responses to questions and that every examiner understands and applies it in the same correct way. As preparation for standardisation each examiner analyses a number of candidates' scripts: alternative answers not already covered by the mark scheme are discussed and legislated for. If, after the standardisation process, examiners encounter unusual answers which have not been raised they are required to refer these to the Principal Examiner.

It must be stressed that a mark scheme is a working document, in many cases further developed and expanded on the basis of candidates' reactions to a particular paper. Assumptions about future mark schemes on the basis of one year's document should be avoided; whilst the guiding principles of assessment remain constant, details will change, depending on the content of a particular examination paper.

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## COMPONENT NUMBER: COMP3

## COMPONENT NAME: Problem Solving, Programming, Operating Systems, Databases and Networking

## STATUS:

Final
DATE:
July 2012
To Examiners:

1. When to award '0' (zero) when inputting marks on CMI+

A mark of 0 should be awarded where a candidate has attempted a question but failed to write anything credit worthy.

Insert a hyphen when a candidate has not attempted a question, so that eventually the Principal Examiner will be able to distinguish between the two (unattempted/nothing credit worthy) in any statistics.
2. This mark scheme contains the correct responses which we believe that candidates are most likely to give. Other valid responses are possible to some questions and should be credited. Examiners should refer off mark scheme responses that they believe are creditworthy to a Team Leader.

The following annotation is used in the mark scheme.
; - means a single mark
// - means alternative response
/ - means an alternative word or sub-phrase
A - means acceptable creditworthy answer
R - means reject answer as not creditworthy
NE - means not enough
I - means ignore
DPT - in some questions a specific error made by a candidate, if repeated, could result in the loss of more than one mark. The DPT label indicates that this mistake should only result in a candidate losing one mark, on the first occasion that the error is made. Provided that the answer remains understandable, subsequent marks should be awarded as if the error was not being repeated.

| 1 |  | Processor management // Allocation of processors // Allocation <br> of processor time // (process) scheduling // thread management; <br> A processing management <br> Allocation/management of RAM / memory // allocation of buffers; <br> Allocation/management of / control of I/O devices/peripherals // <br> I/O management // device driver management; <br> File / backing store / secondary store management / access / <br> organisation; <br> Power / battery management; <br> A Interrupt handling; <br> A Provision of Application Program Interface / API; A interface <br> between hardware and applications <br> A Provision / management of (windows in) user interface; <br> A Management of system security; <br> A Answers by example, only one example of each type <br> A A description of a type of software management but not just <br> "software management". e.g. loading of programs, software <br> installation, registering DLLs. <br> R Software management alone unless role of OS in this is clear <br> e.g. installation of new software, updating registry <br> MAX 3 |  |
| :---: | :--- | :--- | :--- |


| $\mathbf{2}$ | (a) | Backus-Naur (Form); <br> A Backus Normal (Form), BNF, Extended Backus-Naur (Form), <br> Augmented Backus-Naur (Form), ABNF <br> A Misspellings of Backus-Naur <br> A Format for Form and the word "Notation" <br> R BN | 1 |
| :---: | :---: | :--- | :---: |



| 2 | (c) |  | <expression> <br> 1 mark for each area surrounded by a rectangle <br> A missing chevrons <br> DPT Arrows drawn instead of lines | 3 |
| :---: | :---: | :---: | :---: | :---: |


| $\mathbf{3}$ | (a) | Space / Memory (complexity); <br> $\mathbf{A}$ amount of memory used | 1 |
| :---: | :---: | :---: | :--- | :---: |


| 3 | (b) | (i) | \|l| | Pos1 <br> 1 $\square$ $\square$ <br> nswer do no block <br> ddition <br> T if jus unn w 1 is w pe" <br> andida uld be | w1 <br> Rope <br> Dagger $\square$ <br> Rope $\square$ <br> $\square$ <br> which hav award a m which th <br> rows in <br> "Duplicat n it shoul ten in the <br> has not ost (for th | Pos2 <br> 1 <br> 2 <br> 3 <br> 1 <br> 2 <br> 3 <br> 1 <br> 2 <br> 3 | w2  <br> Rope  <br> Dagger  <br> Rope  <br> Rope  <br> Dagger  <br> Rope  <br> Rope  <br> Dagger  <br> Rope  | repeated in empty cells, ny incorrect values within awarded. <br> g as the trace is correct. <br> written in the Output Rope" or if the value of W1 e.g. 1 instead of <br> e of N , only one mark area) for this mistake | 3 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 3 | (b) | (ii) |  |  |  |  |  |  | 1 |
| 3 | (b) | (iii) |  | RKS R PAR ernativ orithm er; A re h loop rnativ (basic ried ou loop rnativ h of th h entry paris pperca nswer 9. <br> heck X 2 | AN ONL <br> 3 (b) (ii) <br> 1: <br> has neste erence to repeats N 2: <br> operatio <br> $\mathrm{N}^{2}$ times; <br> executes <br> 3: <br> ( N ) entri <br> is compa <br> s/operat <br> e or lowe <br> where ex <br> alternativ | BE A <br> loops inner times <br> / If st becau time <br> is c ed N ns ar case mple <br> e to | ARDED <br> / two loo nd outer <br> ement / it is ins <br> mpared nes; so required <br> are used <br> mparison | IF CORRECT ANSWER <br> s with one inside the oops <br> ile read / comparison is de nested loops // because <br> each of the ( N ) others // $N^{2}\left(A N^{*} N\right)$ <br> // $\mathrm{N}^{*} \mathrm{~N}=\mathrm{N}^{2}$; <br> instead of $N$ and $N^{2}$, e.g. 3 | 2 |


| 4 | (a) | Greater the bandwidth, the higher the bit rate // positive <br> correlation // (directly) proportional; <br> Bandwidth must be at least 2wHz where $w$ is the bit rate in bits <br> per second; | MAX <br> 1 |
| :---: | :---: | :---: | :--- | :---: |


| 4 | (b) | Time delay between the moment something is initiated and the <br> moment its effect begins <br> A time delay between signal being transmitted and arriving <br> A time taken for transmitted data to arrive at the receiver <br> A lag for time delay <br> NE delay in transmission, transmission time | 1 |
| :---: | :---: | :---: | :--- | :---: |


| $\mathbf{4}$ | (c) | Bit rate is double/twice baud rate // Baud rate is half bit rate; <br> A "lt" is double; <br> A 2:1 | 1 |
| :---: | :---: | :---: | :--- | :---: |

(d)

| 5 | (a) | (Using an algorithm) to convert a message into a form that is not <br> understandable (without the key to decrypt it); <br> (Using an algorithm) to convert a message into a form that is <br> only understandable by the intended parties // can only be read <br> with the correct key; <br> Converting a message into cipher text; <br> NE scrambling unless further explanation is provided <br> A "unreadable" for "understandable" <br> A "data" for "a message" <br> MAX 1 | 1 |
| :---: | :---: | :--- | :--- |


| 5 | (b) | (i) | B will not be able to decrypt it // A's private key would be needed <br> to decrypt it // only A could decrypt it; (as ...) <br> Only A has access to A's private key // B cannot access A's <br> private key; <br> MAX 1 | 1 |
| :---: | :---: | :---: | :--- | :---: |
|  | (b) | (ii) | As A's public key is available to anyone; <br> Anybody could decrypt it; <br> MAX 1 | 1 |





| 6 | (b) |  | 1 method mark for either: <br> - showing correct value of both mantissa and exponent in denary <br> - showing binary point shifted 6 places to right in binary number <br> - indicating that final answer calculated using answer $=$ mantissa $\times 2^{\text {exponent }}$ <br> Mantissa $=-0.6875 / /-11 / 16$ <br> Exponent $=6$ <br> Answer = -44 <br> 1 mark for correct answer <br> If answer is correct and some working has been shown, award two marks, even if working would not have gained credit on its own. <br> Marks for working can be awarded in the answer. | 2 |
| :---: | :---: | :---: | :---: | :---: |




| 6 | (e) | Calculation |  |  |
| :---: | :---: | :---: | :--- | :--- |
|  |  | Type of <br> Error |  |  |
|  |  | Multiplying two very large numbers together. <br> Overflow; |  |  |
| Dividing a number by a very large number. <br> Adding together two numbers of very different <br> sizes e.g. a tiny number to a very big number. | Underflow; <br> Cancellation; |  |  |  |
| If same answer is used more than once and it is correct in one <br> instance then award the mark for the correct instance. | 3 |  |  |  |


| 7 | (a) |  | Server |
| :--- | :--- | :--- | :--- | :--- |


| 7 | (b) |  | Computer $=$ Class/Subclass/Extends (Device) <br> 1 <br> Accept answers that use different notations, so long as meaning is clear. <br> 1 mark for correct header including name of class and parent class; <br> 1 mark for redefining the AddDevice procedure; <br> 1 mark* for defining all 3 extra functions needed to read variable values, all identified as being public (keyword public is optional if functions are declared before variables); <br> 1 mark $^{\#}$ for defining all 3 extra variables, with appropriate data types and identified as being private; <br> A any sensible numeric types for RAMCapacity and HDDCapacity, do not have to be whole numbers <br> A answers that indicate separately that each variable is private or each method is public <br> R do not award mark for declaring new functions if any of the functions have the same name as the variables <br> I parameters to methods, minor changes to names that do not affect clarity <br> * - Do not award this mark if any extra functions/procedures have been declared, except for functions that would set values e.g. SetProcessorName or an incorrectly named procedure to add e.g. AddComputer <br> \# - Do not award this mark if any extra variables have been declared | 4 |
| :---: | :---: | :---: | :---: | :---: |


| 7 | (c) |  | ```Laptop = Class/Subclass (Computer) (Public) Procedure AddDevice (Override) 1 Function GetBluetoothInstalled Private / Protected 1 \\ BluetoothInstalled : Boolean \\ End``` <br> 1 mark for correct header including name of class and parent class; <br> MAX 1 of the following two marks: <br> 1 mark for redefining the AddDevice procedure; 1 mark* for : <br> - defining the GetBluetoothInstalled function needed to read this value, identified as being public (keyword public is optional if function is declared before variable) <br> - defining the BluetoothInstalled variable with an appropriate data type as being private. <br> A Boolean or whole number types for BluetoothInstalled but reject string, character or real number types <br> A Different sensible name for GetBluetoothInstalled function e.g. CheckBluetoothInstalled, IsBluetoothInstalled <br> A answers that indicate separately that each variable is private or each method is public <br> I parameters to methods, minor changes to names that do not affect clarity <br> I addition of any extra functions or variables <br> * Do not award this mark if any extra functions / procedures / variables declared, except for a SetBluetoothInstalled procedure. | 2 |
| :---: | :---: | :---: | :---: | :---: |


| 7 | (d) | What (2 marks): <br> Wireless/RF (protocol/standard/technology); <br> For exchanging data over short distances // for creating <br> Personal Area Network; <br> NE "uses waves" for "wireless" |  |
| :--- | :--- | :--- | :--- |
| Example (1 mark): <br> Any sensible example, related to the use of Bluetooth with the <br> laptop e.g. synchronising contacts between phone/laptop, <br> sending photographs from phone to laptop, Bluetooth mouse, <br> Bluetooth headset/headphones (used with laptop) etc; NE <br> connecting to wireless network NE mouse <br> If the example makes clear that the technology is wireless, but <br> this is not explicitly stated in the "What" part of the response <br> then the "Wireless" mark should be awarded in the "What" part. | 3 |  |  |


| 8 | (a) | System will be storing confidential/personal data (that must be <br> kept securely/safely); <br> Centralised/improved security management // centralised login <br> system // centralised administration // administration will be <br> easier; <br> Centralised backup; <br> Harder for users to change security/sharing settings; <br> Running database from a server will avoid concurrency issues // <br> will avoid problems if two users/computers update (a record in <br> the) database simultaneously; A will allow simultaneous <br> updates/access R answers that imply that on a peer-to-peer <br> system there would be a separate copy of the database on each <br> workstation <br> Running database from server will ensure that it is always <br> available (as server is unlikely to be turned off) // Files would <br> always be available (as server is unlikely to be turned off); <br> Server (operating system) may allow more simultaneous <br> connections than a workstation // (operating system software on) <br> workstations may not allow enough simultaneous connections <br> for ten users; |  |  |
| :---: | :---: | :--- | :--- | :--- |
| NE the database could be stored on the server |  |  |  |  |
| MAX 2 | 2 |  |  |  |




| $\mathbf{8}$ | (c) | To connect networks using different protocols // to convert <br> transmitted data from one protocol to another; | 1 |
| :---: | :---: | :---: | :--- | :---: |


| 9 | (a) | Composite (key); A Compound (key) <br> Note: The word key is not required | 1 |
| :---: | :---: | :--- | :--- | :--- |


| 9 | (b) | Data is atomic // no repeating groups (of attributes); R No <br> repeated columns/attributes/data/values <br> No partial (key) dependencies // No (non-key) attribute depends <br> on part of the primary key but not the whole of it // all non-prime <br> attributes are (functionally) dependent on the whole of every <br> candidate key // (non-key) attributes depend on the whole key; <br> No non-key dependencies // No transitive dependencies // (non- <br> key) attributes depend on nothing but the key; <br> Every (non-key) attribute is dependent upon the key; <br> Every determinant is a candidate key; <br> A "field" for "attribute" <br> A "part" for "partial" <br> MAX 2 |  |
| :---: | :---: | :--- | :--- |


| 9 | (c) |  | 1 mark per correct relationship (the dashed one is given) <br> MAX 1 if more than two relationships drawn | 2 |
| :---: | :---: | :---: | :---: | :---: |


| 9 | (d) |  | Solution 1: <br> SELECT EmailAddress, Forename, Surname <br> FROM Book, Member, Loan <br> WHERE Author = 'Lucas Bailey' AND <br> Book.BookID=Loan.BookID AND <br> Member.MemberID=Loan.MemberID <br> 1 mark for correct three fields in SELECT clause <br> 1 mark for correct three tables in FROM clause <br> 1 mark for Author = 'Lucas Bailey' <br> 1 mark for Book.BookID=Loan.BookID linked by AND <br> 1 mark for Member.MemberID=Loan.MemberID linked by AND <br> Solution 2: <br> SELECT EmailAddress, Forename, Surname <br> FROM Book INNER JOIN Loan ON Book.BookID=Loan.BookID <br> INNER JOIN Member on <br> Member.MemberID=Loan.MemberID <br> WHERE Author = 'Lucas Bailey' <br> 1 mark for correct three fields in SELECT clause <br> 1 mark for correct three tables in FROM clause <br> 1 mark for join from Member to Loan <br> 1 mark for join from Loan to Book <br> 1 mark for Author = 'Lucas Bailey' <br> Note: Joins do not need to be done in same order as example <br> Do not award mark for SELECT clause if extra attributes listed. <br> Do not award mark for 'Lucas Bailey' unless it is enclosed in single or double quotation marks. <br> Accept table names before fieldnames. <br> Accept use of Alias/AS command e.g. FROM Member as M then use of M as table name. <br> Accept insertion of spaces into fieldnames <br> DPT for unnecessary punctuation - allow one semicolon at the very end of the statement, but not at the end of each clause. <br> Also, allow insertion of brackets at logically allowable places in the WHERE/FROM clauses. <br> DPT for fieldname before table name. <br> Refer responses using nested SQL queries to team leaders. | 5 |
| :---: | :---: | :---: | :---: | :---: |


| 9 | (e) | Alternative 1: <br> INSERT INTO Book <br> VALUES ( 837023, "Kenyan Safari", "Karen Matu", "African <br> Travel Guides" ) <br> Alternative 2: <br> INSERT INTO Book (BookID, Title, Author, Publisher) <br> VALUES ( 837023, "Kenyan Safari", "Karen Matu", "African <br> Travel Guides" ) <br> 1 mark for INSERT INTO Book; <br> 1 mark for correct field values. If alternative 2 is used, the order <br> of the values and fieldnames must correspond to each other; |  |
| :--- | :--- | :--- | :--- |
| The values Kenyan Safari, Karen Matu and African Travel <br> Guides must be in single or double quotation marks for the mark <br> to be awarded. Accept the value 837023 with or without <br> quotation marks. <br> A Minor errors in transcribing the data from the question into the <br> answer. <br> A omission of brackets |  |  |  |


| (f) | ONE MARK FOR PRINCIPLE AND MAX TWO MARKS FOR <br> IMPLEMENTATION <br> Principle: <br> Create a new table (A link table) (BookCopy); through which <br> Book and Loan tables will be (indirectly) linked; <br> Implementation details using a new primary key: <br> Create a new unique ID/key field (e.g. CopyID) (for each copy); <br> Store the BookID and the CopyID in the new table; <br> Replace the BookID in the Loans table with this CopyID; <br> Note: In this implementation, CopyID is unique, i.e. BookID 1 <br> and 2 cannot both have CopyID 1. <br> Implementation details using a composite key: <br> Create a new field CopyID; <br> Composite key formed by BookID and CopyID; TO if composite <br> key is clearly in book table or loan table <br> Store the BookID and the CopyID in the new table; R adding <br> CopyID to Book table as this would created data redundancy but <br> this does not talk out other mark scheme points <br> Add the CopyID field to Loans table; R replace BookID with <br> CopyID <br> Note: In this implementation, CopyID is not unique, e.g. BookID <br> 1 and 2 can both have CopyID 1. |  |
| :---: | :---: | :--- | :--- |


|  | MAX 3 <br> Marks can be awarded for principle and/or implementation <br> details. |  |
| :--- | :--- | :--- | :---: |
| A Relation for Table <br> A Answers if candidates have rewritten new relations, awarding <br> marks where the points above can be observed in the redrawn <br> relations; <br> A alternative name for CopyID | 3 |  |


| $\mathbf{9}$ | (g) | (i) | So that searching, adding and deleting can be done efficiently // <br> To speed up searching, adding and deleting; <br> A just one of searching, adding, deleting <br> NE organise efficiently <br> NE easily for efficiently |  |
| :--- | :--- | :--- | :--- | :--- |
| (g) | (ii) | Alternative 1 (context-specific): <br> A function/calculation that computes a record position/address; <br> within a specified range; from a key field value; <br> A an example of a hashing function e.g. calculate an integer <br> from certain letters in a field for one mark |  |  |
| (g) | Alternative 2 (generic): <br> A function (A algorithm) H, applied to a key k; which generates a <br> hash value (H(k)) (of range smaller than the domain of values of <br> k); <br> (iii) | MAX 2 | What is (1 mark): <br> When more than one key value maps to the same record <br> position/address // when two keys compute the same hash <br> value; <br> A "two records", "two items" or "two pieces of data" for "two <br> keys" but R "two files" - both in this question part only | 2 |
| How dealt with (1 mark): <br> Store the record in the next available location in the file // store a <br> pointer (in each file location) that points to a list of records that have <br> all collided at the file location; <br> A idea that each storage location could store more than one record <br> e.g. five records per location, if explained. <br> A example of what "next available" might be <br> A key is rehashed | A table for file |  |  |  |


| 10 | (a) | Connected // There is a path between each pair of vertices; <br> Undirected // No direction is associated with each edge; <br> Has no cycles // No (simple) circuits // No closed chains // No <br> closed paths in which all the edges are different and all the <br> intermediate vertices are different // No route from a vertex back <br> to itself that doesn't use an edge more than once or visit an <br> intermediate vertex more than once; A no loops <br> MAX 1 <br> Alternative definitions: <br> A simple cycle is formed if any edge is added to graph; <br> Any two vertices can be connected by a unique simple path; | 1 |
| :---: | :---: | :---: | :--- | :---: |


| $\mathbf{1 0}$ | (b) | No route from entrance to exit / through maze; <br> Maze contains a loop/circuit ; A more than one route through <br> maze; <br> Part of the maze is inaccessible / enclosed; <br> R Responses that clearly relate to a graph rather than the maze <br> MAX 1 | 1 |
| :---: | :---: | :---: | :--- | :---: |



| $\mathbf{1 0}$ | (d) | (i) | Routine defined in terms of itself // Routine that calls itself; <br> A alternative names for routine e.g. procedure, algorithm <br> NE repeats itself | 1 |
| :---: | :---: | :---: | :--- | :---: |
| $\mathbf{1 0}$ | (d) | (ii) | Stores return addresses; <br> Stores parameters; <br> Stores local variables; NE temporary variables <br> Stores contents of registers; <br> A To keep track of calls to subroutines/methods etc. <br> MAX 1 <br> Procedures / invocations / calls must be returned to in reverse <br> order (of being called); <br> As it is a LIFO structure; A FILO <br> As more than one / many return addresses / sets of values may <br> need to be stored (at same time) // As the routine calls itself and <br> for each call/invocation a new return address / new values must <br> be stored; <br> MAX 1 | 2 |



| 11 | (a) | Is it possible in general to write a program/algorithm; that can <br> tell, given any program and its inputs and without <br> running/executing the program;, whether the given program <br> with its given inputs will halt? | A "it" in second reference to program. <br> A "create a Turing machine" for "write an algorithm" | 2 |
| :---: | :---: | :---: | :--- | :---: |


| 11 | (b) | Shows that some problems are non-computable / undecideable <br> // shows that some problems cannot be solved by a <br> computer/algorithm; <br> In general, inspection alone cannot always determine whether <br> any given algorithm will halt for its given inputs // a program <br> cannot be written that can determine whether any given <br> algorithm will halt for its given inputs; <br> A it is not computable <br> MAX 1 | 1 |
| :---: | :---: | :---: | :--- | :---: |


| $\mathbf{1 2}$ | (a) | $a^{+} b^{+} / / a b b^{*} \mathrm{c} / / \mathrm{ab}{ }^{*} \mathrm{bc} ;$ <br> $\mathbf{I}^{\wedge}$ at start, $\$$ at end of expression | 1 |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{1 2}$ | (b) | $(0 \mid 1) 1^{*} / /(1 \mid 0) 1^{*} / /[01] 1^{*} / /[10] 1^{*} / /[0 \mid 1] 1^{*} / /[1 \mid 0] 1^{*} / / 0 \mid\left(0 ? 1^{+}\right)$ <br> $\mathbf{l}^{\wedge}$ at start, $\$$ at end of expression | 1 |

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